## GENERALIZED CONVEX KERNELS

#### BY

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### ABSTRACT

The notion of the convex kernel of a set D is generalized to that of the *n*-th order kernel of D. Such kernels are studied for compact, simply connected subsets of the Euclidean plane. In particular, it is shown that under certain circumstances [see Theorem 4 and also section 5], these kernels have rather simple structures.

1. Introduction. Horn and Valentine [3] have generalized the notion of convex set to that of  $L_n$  set. A set D in the Euclidean plane  $E_2$  is called an  $L_n$  set if for every pair of points x and y in D, there is a polygonal line of at most n segments lying in D which joins x to y. Such sets can be used to approximate (in the sense of the Hausdorff metric) any compact, connected set (see [1]). The results of [1] have been extended by McCoy [4] to complete locally compact convex metric spaces. For some work concerning partitions of Euclidean spaces into  $L_n$  sets, the reader is referred to Ceder [2].

The convex kernel of a set D is defined to be the set of points x in D such that for all y in D, the segment joining x to y is contained in D. It is well known that the convex kernel of a set is itself convex.

The purpose of this note is to study the *n*th order kernel of D, by which we mean the set of points x in D such that for each y in D there is a polygonal line of at most n segments lying in D, and joining x to y.

If D is the boundary of a square, for example, the 2nd order kernel consists of the four corners of the square. By considering examples similar to this, one can easily verify that the *n*th order kernel of a set need not even be connected if n > 1. The essential difference between the cases n = 1 and n > 1 is that there is a unique line determined by any pair of points x and y, but there is an infinitude of polygonal lines with at most n segments joining x to y if n > 1. As we shall see, the assumption that D is simply connected partially overcomes this difficulty.

2. Notation and terminology. In the sequel, D will denote a compact, simply connected set in  $E_2$ . If B is a set, then  $\delta B$  will denote its boundary and  $\sim B$  its complement. We shall use the notation  $\langle p_0, p_1, ..., p_n \rangle$  for the *n*-sided polygonal

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line (n-line) joining  $p_0$  to  $p_n$  with  $p_1, ..., p_{n-1}$  consecutive, intermediate vertices It will always be assumed that  $\langle p_0, ..., p_n \rangle$  has no self intersections. If x and y are points in D, then by  $\rho(x, y)$  we shall mean the minimum number of segments that a polygonal line lying in D and joining x to y can have. A full  $L_n$  set is an  $L_n$  set for which there exist  $x, y \in D$  with  $\rho(x, y) = n$ . If  $x \in D$ , then  $K_x^n = \{y: \rho(x, y) \leq n\}$ . If no confusion can arise, we shall write  $K_x$  in place of  $K_x^n$ . When considering a sequence  $\{x_i\}$  of points of D, we shall, for simplicity of notation, write  $K_i$  for  $K_{x_i}$ . We shall denote by  $K^n$  the nth order kernel of D. Clearly  $K^n = \bigcap_{x \in D} K_x^n$ .

3. Some properties of  $K_x$ . In this section we record four lemmas which we shall use in the next section. Lemma 1, below, will be used frequently in the sequel.

LEMMA 1. If  $x \in D$  and  $y, z \in K_x$  and  $\langle y, z \rangle \subset D$ , then  $\langle y, z \rangle \subset K_x$ .

**Proof.** Let  $L_y = \langle x, y_1, ..., y_{n-1}, y \rangle$  and  $L_z = \langle x, z_1, ..., z_{n-1}, z \rangle$  be *n*-lines in D joining x to y and z respectively. The n-lines  $L_y$  and  $L_z$  along with  $\langle y, z \rangle$  determine a figure P which is the union of a finite number of simple closed polygons (with interiors) some of which may degenerate into segments. Since D is simply connected,  $P \subset D$ . If  $z \in \langle y_{n-1}, y \rangle$  or  $y \in \langle z_{n-1}, z \rangle$ , the conclusion follows trivially. Likewise the conclusion is immediate if either  $L_y$  or  $L_z$  intersects  $\langle y, z \rangle$  other than at the endpoints. If not, then let  $w \in \langle y, z \rangle$ . The point w is in one of the polygons P' of P. Each of the vertices  $u_1, ..., u_r$  of P' is one of the  $y_i$ , one of the  $z_j$ or a point at which a segment of  $L_y$  intersects a segment of  $L_z$ . Let  $T = \{v \in P' : \langle w, v \rangle \subset P'\}$ . It is clear that there exists  $q, 1 \leq q \leq r$ , such that  $u_q \in T, u_q \neq y, u_q \neq z$ . If for some s,  $1 \leq s \leq n-1, u_q = y_s$  then the (s + 1)-line  $\langle x, y_1, ..., y_s, w \rangle$  is contained in P and  $w \in K_x$ . A similar n-line exists in case  $u_q = z_s$  for some s. If  $u_q$  is a point of intersection of  $L_y$  and  $L_z$ , then the extension of the segment  $\langle u_q, w \rangle$  (into another polygon P" of P) intersects the boundary of P at a point t. It is easy to verify that there is an n-line lying in P having terminal side  $\langle t, w \rangle$ .

# LEMMA 2. If $x \in D$ , then $K_x^n$ is a compact, simply connected $L_{2n}$ set.

**Proof.** The compactness of  $K_x^n$  is obvious.

That  $K_x$  is an  $L_{2n}$  set follows trivially from the fact that any two points of  $K_x^n$  can be joined by a 2n-line whose middle vertex is x.

It remains to show that  $\sim K_x^n$  has no bounded components. Assume  $\sim K_x^n$  has a bounded component C. Let  $y \in C$  and let  $\langle y_1, y_2 \rangle$  be a segment through y where  $y_1, y_2 \in \delta C$ , and there are no other points of  $\delta C$  on  $\langle y_1, y \rangle$  and  $\langle y, y_2 \rangle$ . Since boundary points of C must also be boundary points of  $K_x^n$ , and since  $K_x^n$  is closed, it follows that  $y_1$  and  $y_2$  are in  $K_x^n$ . Now each point z of  $\langle y_1, y_2 \rangle$  is in D, otherwise the component of  $\sim D$  containing z would be bounded. By Lemma 1,  $y \in K_x^n$  contradicting the hypothesis  $y \in C$ . Hence  $K_x^n$  is simply connected. LEMMA 3. Let  $x_1, x_2, ..., x_m$  be points of D and let  $M = K_1 \cap K_2 \cap ... \cap K_{m-1}$ . Then  $\sim (M \cup K_m)$  has no bounded components.

**Proof.** Suppose that  $\sim (M \cup K_m)$  has a bounded component C. Since C is a bounded, open, connected set, it can be shown that there exist three points  $y_1, y_2, y_3 \in \delta C$  such that the open segments  $(y_1, y_2), (y_2, y_3)$  and  $(y_1, y_3)$  are contained in C. Now since  $\delta C \subset M \cup K_m$ , two of these three points, say  $y_1$  and  $y_2$ , are in M or in  $K_m$ . In either case, it follows from Lemma 1 that the entire segment  $\langle y_1, y_2 \rangle$  is in M or in  $K_m$ . This contradicts the existence of a bounded component of  $\sim (M \cup K_m)$ .

LEMMA 4. Let  $x_1, \ldots, x_m$  be points of D. Then  $K_1^n \cap K_2^n \cap \cdots K_m^n$  is a compact, simply connected  $L_{2n}$  set.

**Proof.** That  $K_1^n \cap ... \cap K_m^n$  is compact follows trivially from the compactness of  $K_1^n,...,K_m^n$ . If m = 1, the theorem reduces to Lemma 2. Let  $M = K_1^n \cap ... \cap K_{m-1}^n$ . Assume M is a simply connected,  $L_{2n}$  set. Let y, z be points of  $M \cap K_m^n$ ; let  $L = \langle y, y_{n-1}, ..., y_1, x_m, z_1, ..., z_{n-1}, z \rangle$  and  $L_M = \langle y, v_1, v_2, ..., v_{2n-1}, z \rangle$  be 2*n*-lines joining y and z lying in  $K_m^n$  and M respectively. These lines determine a figure Pwhich is a union of a finite number of simple closed polygons with interiors some of which may degenerate into segments. We may further assume that  $L_M$  is simple and no line segment joining nonadjacent vertices of  $L_M$  lies entirely in P. Since  $\sim (M \cup K_m^n)$  has no bounded components,  $P \subset M \cup K_m^n$ . We shall show that there is a polygonal path in  $K_m^n \cap M$  having at most the number of segments of  $L_M$ .

For each  $v \in L_M$  let

$$s(v) = \sup \{t \in L_M : \langle v, t \rangle \subset P \text{ and } (v, t) \cap \langle z_{n-1}, z \rangle \text{ is empty} \}$$

where the supremum is taken with respect to the natural ordering of  $L_M$  from y to z. We shall denote this ordering by " $\prec$ ".

We first show that if  $v \in K_m$ , then  $s(v) \in K_m^n$ . Now, if  $s(v) \in \langle z_{n-1}, z \rangle$ , the conclusion is trivially true. If  $s(v) \notin \langle z_{n-1}, z \rangle$ , then for some point p on L, the points v s(v) and p are collinear. If  $p \in \langle z_{n-1}, z \rangle$ , then  $s(v) \in \langle v, p \rangle$  and by Lemma 1,  $s(v) \in K_m^n$ . If  $p \in \langle z_m, z_{m+1} \rangle$ ,  $m \neq n-1$ , then the polygonal line  $\langle x_m, z_1, ..., z_m, p, s(v) \rangle$  lies in D and has at most n segments so that  $s(v) \in K_m^n$ . A similar polygonal path exists if  $p \in \langle y_{m+1}, y_m \rangle$   $(m \neq n-1$  unless  $p = y_{n-1})$ .

It is clear that if  $v_i \leq v < v_{i+1}$  and  $(v_i, v_{i+1}) \cap \langle z_{n-1}, z \rangle$  is empty, then  $v_{i+1} \leq s(v)$ . If  $s(v) \in \langle v_i, v_{i+1} \rangle \cap \langle z_{n-1}, z \rangle$ , then the line  $\langle v, s(v), z \rangle$  has at most as many segments (two or, in the degenerate case s(v) = z, one) as  $\langle v, v_{i+1}, ..., v_{2n-1}, z \rangle$ . Thus the polygonal line

$$\langle y, s(y), s^2(y), ..., z \rangle$$

has at most 2n segments. It follows from Lemma 1 that this 2n-line lies in  $K_m^n \cap M$ .

The simple connectedness of  $K_m^n \cap M$  now follows from the simple connectedness of  $K_1, \ldots, K_m$ .

4. The *n*th order kernel. We now proceed to study the *n*th order kernel of D. In Theorems 1 and 3 we obtain representations for  $K^n$ . Theorem 2 shows that  $K^n$  shares some of the properties of the sets  $K_x$  obtained in Lemma 2. In Theorem 4 we show that for a full  $L_{2n}$  set,  $K^n$  is itself an  $L_2$  set.

THEOREM 1. Let D' be a dense subset of D. Then  $K^n = \bigcap_{x \in D'} K_x^n$ .

**Proof.** Let  $R = \bigcap_{x \in D} K_x^n$ . Trivially,  $K^n \subset R$ . Now assume  $y \in R$ . Then  $y \in K_x^n$  for each  $x \in D'$ . Equivalently  $D' \subset K_y^n$ . But  $K_y^n$  is closed so that  $K_y^n$  contains the closure of D' which is D. Thus  $D \subset K_y^n$  and  $y \in K^n$  so that  $R \subset K^n$ .

THEOREM 2. The nth order kernel is a compact, simply xconnected,  $L_{2n}$  set.

**Proof.** Let  $D' = \{x_1, ..., x_m, ...\}$  be a countable, dense subset of D; let  $M_m = K_1^n \cap ... \cap K_m^n$ . The sets  $M_m$  form a decreasing sequence of compact,  $L_{2n}$  sets whose intersection  $K^n$  is therefore a compact,  $L_{2n}$  set (see [1]: Theorem 2). The simple connectedness of K follows immediately from the connectedness of  $K^n$  and the simple connectedness of  $K_1^n, K_2^n, ...$ 

THEOREM 3. Let  $S = \bigcap K_x^n$ , the intersection being taken over all points  $x \in \delta D$ . Then  $S = K^n$ .

**Proof.** Trivially  $K^n \subset S$ . Let  $x \in D \sim K^n$  and let  $y \in \sim K_x^n$ . If  $y \in \delta D$  then  $x \in D \sim S$ . If y is an interior point of D, then the component of  $\sim K_x^n$  containing y must contain a boundary point of D, for otherwise  $K_x^n$  would not be simply connected. Thus  $x \notin S$  and  $S \subset K^n$ .

The corollary below generalizes [3: Theorem 1.4].

COROLLARY. If  $y \in K_x^n$  for all x,  $y \in \delta D$  then D is an  $L_n$  set.

**Proof.** If, for every  $x \in \delta D$ ,  $y \in K_x^n$  for all  $y \in \delta D$ , then by Theorem 3,  $\delta D \subset K^n$ Since  $K^n$  is simply connected, it follows that  $D \subset K^n$ . Thus for each  $x, y \in D$  we have  $\rho(x, y) \leq n$ .

LEMMA 5. Let D be a full  $L_{2n}$  set, let x,  $y \in K^n$  and let  $\alpha$  be a point of D such that  $\rho(\alpha, x) = \rho(\alpha, y) = n$ . Then  $K_x^1 \cap K_y^1 \cap K_\alpha^n$  is non empty.

**Proof.** We first show that  $K_x^1 \cap K_y^1$  is non empty. Since *D* is a full  $L_{2n}$  set, there are points  $\beta$  and  $\gamma$  with  $\rho(\beta,\gamma) = 2n$ . Let  $L_y = \langle \beta, b_{n-1}, ..., b_1, y, c_1, ..., c_{n-1}, \gamma \rangle$  and  $L_x = \langle \beta, \beta_{n-1}, ..., \beta_1, x, \gamma_1, ..., \gamma_{n-1}, \gamma \rangle$  be 2*n*-lines joining  $\beta$  and  $\gamma$  via *x* and *y* respectively. The lines  $L_x$  and  $L_y$  determine a figure *P* which is the union of a finite number of simple closed polygons with interiors, some of which may degenerate into segments.

If  $x \in L_y$  (or  $y \in L_x$ ) then  $\langle x, y \rangle \subset L_y(\langle x, y \rangle \subset L_x)$ , for otherwise we would have  $\rho(\beta, \gamma) < 2n$ ; hence  $\langle x, y \rangle \subset K_x^1 \cap K_y^1$ .

If  $x \notin L_y$  and  $y \notin L_x$  then x and y are accessible to the interior of P. There must be a segment  $\langle x, v \rangle$  lying in P with  $v \in L_y$  since  $\rho(\beta, \gamma) = 2n$ . Similarly there is a segment  $\langle y, t \rangle \subset P$  with  $t \in L_x$ . It is immediately clear that

and  
$$t \in \langle \beta_2, \beta_1, x, \gamma_1, \gamma_2 \rangle$$
$$v \in \langle b_2, b_1, y, c_1, c_2 \rangle.$$

Several situations can occur.

If  $t \in \langle \beta_1, x, \gamma_1 \rangle$  (or  $v \in \langle b_1, y, c_1 \rangle$ ) then  $v \in K_x^1 \cap K_y^1$ ) ( $t \in K_x^1 \cap K_y^1$ ).

If  $t \in \langle \beta_2, \beta_1 \rangle$  and  $v \in \langle b_2, b_1 \rangle$  (or  $t \in \langle \gamma_1, \gamma_2 \rangle$  and  $v \in \langle c_1, c_2 \rangle$ ) then it can be verified that  $\langle x, v \rangle \cap \langle y, t \rangle = \{p\} \in P$  and  $p \in K_x^1 \cap K_y^1$ .

If  $t \in \langle \beta_2, \beta_1 \rangle$  and  $v \in \langle c_1, c_2 \rangle$  (or  $t \in \langle \gamma_1, \gamma_2 \rangle$  and  $v \in \langle b_2, b_1 \rangle$ ) then it can be verified that one of the vertices of the hexagon  $\langle y, t, \beta_1, x, v, c_1, y \rangle$  ( $\langle y, t, \gamma_1, x, v, b_1, y \rangle$ ) is in  $K_x^1 \cap K_y^1$ .

We proceed to show that  $K_x^1 \cap K_y^1 \cap K_\alpha^n$  is non empty. Let p be a point of  $K_x^1 \cap K_y^1$ . The lines  $L_1 = \langle x, a_{n-1}, ..., a_1, \alpha, \alpha_1, ..., \alpha_{n-1}, y \rangle$  and  $L_2 = \langle x, p, y \rangle$  determine a figure  $P^*$ . An analysis of  $P^*$  similar to the above analysis of P, but using the fact that any polygonal path between x and  $\alpha$  or y and  $\alpha$  must have at least n segments, verifies that there is a point  $v \in P^*$  which lies in  $K_x^1 \cap K_y^1 \cap K_n^n$ .

**THEOREM 4.** Let D be a full  $L_{2n}$  set. Then  $K^n$  is an  $L_2$  set.

**Proof.** If  $K^n$  is not an  $L_2$  set, then there are points  $x, y \in K^n$  such that for each  $t \in K_x^1 \cap K_y^1$ , there exists  $\alpha(t) \in D$  with  $t \in \sim K_{\alpha(t)}^n$ . Since  $K_x^1 \cap K_y^1$  is compact, there exist points  $\alpha_1, \alpha_2, ..., \alpha_p \in D$  such that for each  $t \in K_x^1 \cap K_y^1$ ,  $t \in \sim K_i^n$  for some  $i, 1 \leq i \leq p$ . Each of the  $\alpha_i$  satisfies  $\rho(\alpha_i, x) = \rho(\alpha_i, y) = n$ .

We will show that  $K_1^n \cap K_2^n \cap ... \cap K_q^n \cap K_x^1 \cap K_y^1$  is not empty for q = 1, 2, ...Now it follows from Lemma 5, that for each  $\alpha_i$  there is a point  $t_i \in K_x^1 \cap K_y^1 \cap K_i^n$ . Hence the conclusion is valid for q = 1. Assume the conclusion holds for q - 1. Then there exists  $t \in K_1^n \cap ... \cap K_{q-1}^n \cap K_x^1 \cap K_y^1$ . By Lemma 5 we have the existence of t' in  $K_q^n \cap K_x^1 \cap K_y^1$ . The 2-lines  $L = \langle x, t, y \rangle$  and  $L' = \langle x, t', y \rangle$ determine a figure P which is the union of at most two simple closed polygons with interior, some of which may degenarete into segments. Since D is simply connected,  $P \subset D$ . Now  $\langle x, \rangle$  is not contained in P. For then we would have  $\langle x, y \rangle \subset K^n$  contrary to the assumption that x and y cannot be joined by a 2-line in  $K^n$ .

If L and L' intersect at a point  $v, v \neq x, y$ , then  $v \in K_1^n \cap ... \cap K_q^n \cap K_x^1 \cap K_y^1$ .

If  $L \cap L' = \{x\} \cup \{y\}$  then for one of the lines L (or L') we have  $\langle x, t(\text{or } t'), y \rangle \sim \{x\} \sim \{y\} \subset \operatorname{Int} P^*$  where  $P^*$  is the figure determined by  $\langle x, t'(\text{or } t), y \rangle$  and  $\langle x, y \rangle$ . It is easily verified that  $t(\operatorname{or } t') \in K_1^n \cap \ldots \cap K_q^n \cap K_x^1 \cap K_y^1$ . By induction there is a point  $v \in K_1^n \cap \ldots \cap K_p^n \cap K_x^1 \cap K_y^1$  contrary to the assumption that the sets  $\sim K_i^n \operatorname{cover} K_x^1 \cap K_y^1$ .

5. Kernels of nowhere dense sets. Throughout this section D is nowhere dense

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in addition to being compact and simply connected. The following result is easy to verify and the proof is ommitted.

LEMMA 6. Let x,  $y \in D$ . If  $L_p$  and  $L_q$  are p- and q-lines respectively,  $p \leq q$ , joining x to y and  $L = \{t: t \in L_p \cap L_q\}$  then L is an r-line joining x to y with  $r \leq p$ .

The following theorems can be proved by applications of Lemma 6.

THEOREM 5. Let D be a full  $L_{2n} [L_{2n-1}]$  set; let x,  $y \in D$  with  $\rho(x, y) = 2_n [\rho(x, y) = 2n - 1]$ . If  $\langle x, \alpha_1, ..., \alpha_n, ..., \alpha_{2n-1}, y \rangle [\langle x, \alpha_1, ..., \alpha_{n-1}, \alpha_n, ..., \alpha_{2n-2}, y \rangle$  is a 2n-line [(2n - 1)-line] joining x to y then  $\alpha_n \in K^n [\langle \alpha_{n-1}, \alpha_n \rangle \subset K^n]$ .

THEOREM 6. Let D be a full  $L_{2n}[L_{2n-1}]$  set. Then  $K^n$  is a single point  $[K^n]$  is a single segment].

THEOREM 7. Let D be a full  $L_{2n}[L_{2n-1}]$  set and for p > n, let  $K^p$  denote its pth order kernel; let p = n + q. Then  $K^p$  is an  $L_{2q}[L_{2q+1}]$  set.

A basic difference between the cases in which D is nowhere dense and the general case is that in the former, any two points of D determine a unique path of fewest segments (as Lemma 6 illustrates), whereas this is not so in the general case. Theorems 5 and 6 obviously have no counterparts in the general case. We suspect that Theorem 7 does have an analogue in the general case but have been unable to prove this. It is worth noting that Theorem 6 implies that the *n*th order kernel of a full  $L_{2n}[L_{2n-1}]$  set is non empty in case D is nowhere dense. This is not necessarily true in the general case. For example, if D is the simply connected set determined by a triangle whose sides are extended one unit in each direction, then D is a full  $L_2$  set with an empty first order (convex) kernel.

6. Some conclusing remarks. We conclude with several observations. Simple examples show that  $K^n$  might be contained entirely in the interior of D or entirely in  $\delta D$ , even if D is bounded by a simple Jordan curve.

It can be shown that if  $x \in D$  then the boundary of a component of  $D \sim K_x^n$  can be decomposed into two sets A and B, where  $A \subset \delta D$  and B is either a subinterval of a single segment of an *n*-line in  $K_x^n$  or empty. No corresponding statement can be made for the boundary of a component of  $D \sim K^n$ .

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